#### LECTURE 21

# Scheduling in Distributed Systems



#### General

- Scheduling refers to assigning a resource and a start time end to a task
- Much of scheduling research has been done in Operations Research Community e.g Job Shop, Flow shop scheduling etc.
- Scheduling is often an overloaded term in Grids.
- A related term is mapping that assigns a resource to a task but not the start time.

#### Systems taxonomy

- Parallel Systems
- Distributed Systems
- Dedicated Systems
- Shared Systems
  - Time Shared e.g. aludra
  - Space Shared e.g. HPCC cluster
- Homogeneous Systems
- Heterogeneous Systems

## **Scheduling Regimens**

- Online/Dynamic Scheduling
- Offline/Static Scheduling
- Resource level Scheduling
- Application level Scheduling

## **Applications taxonomy**

- Bag of tasks Independent tasks
- Workflows dependent tasks
  - Generally Directed Acyclic Graphs (DAGs)
- Performance criteria
  - Completion time (makespan), reliability etc.

#### Scheduling Bag of Tasks on Dedicated Systems

- Min-Min
- Max-Min
- Sufferage



#### Min-Min Heuristic

- For each task determine its minimum completion time over all machines
- Over all tasks find the minimum completion time
- Assign the task to the machine that gives this completion time
- Iterate till all the tasks are scheduled

## **Example of Min-Min**

	Τ1	T2	Т3
M1	140	20	60
M2	100	100	70

	T1	Т3
M1	160	80
M2	100	70

	T1
M1	160
M2	170

Stage 1:

- T1-M2 = 100
- T2-M1 = 20

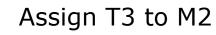
Stage 2: T1-M2 = 100

T3-M2 = 70

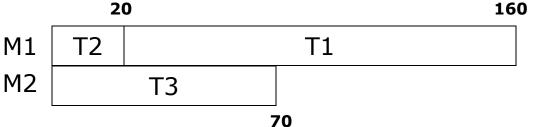
Stage 3: T1-M1 = 160

T3-M1 = 60

Assign T2 to M1



Assign T1 to M1





### Max-Min Heuristic

- For each task determine its minimum completion time over all machines
- Over all tasks find the maximum completion time
- Assign the task to the machine that gives this completion time
- Iterate till all the tasks are scheduled

## **Example of Max-Min**

	Τ1	T2	Т3
M1	140	20	60
M2	100	100	70

	T2	Т3
M1	20	60
M2	200	170

	T2
M1	80
M2	200

Stage 1: T1-M2 = 100

T2-M1 = 20

T3-M1 = 60

Assign T1 to M2

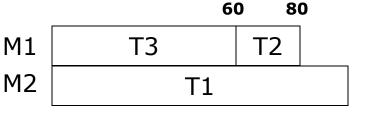
Stage 2: T2-M1 = 20

T3-M1 = 60

Assign T3 to M1

Stage 3: T2-M1 = 80

Assign T2 to M1



<sup>100</sup> 

## Sufferage Heuristic

- For each task determine the difference between its minimum and second minimum completion time over all machines (sufferage)
- Over all tasks find the maximum sufferage
- Assign the task to the machine that gives this sufferage
- Iterate till all the tasks are scheduled

## Example of Sufferage

	Τ1	T2	Т3
M1	140	20	60
M2	100	100	70

	T1	Т3
M1	160	80
M2	100	70

Stage 2:

T1 = 60

T3 = 10

	Т3
M1	80
M2	170

Stage 3:

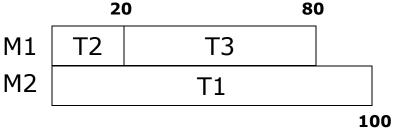
T3 = 90

Stage 1:

- T1 = 40
- T2 = 80

T3 = 10





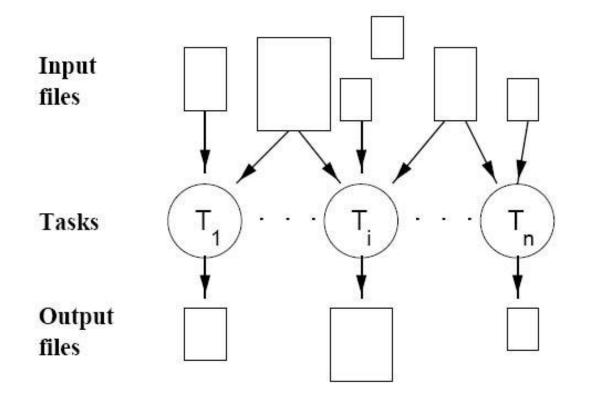
Assign T3 to M1

## **Grid Environments**

- Time-shared resources
- Heterogeneous resources
- Tasks require input files that might be shared
- Data transfer times are important

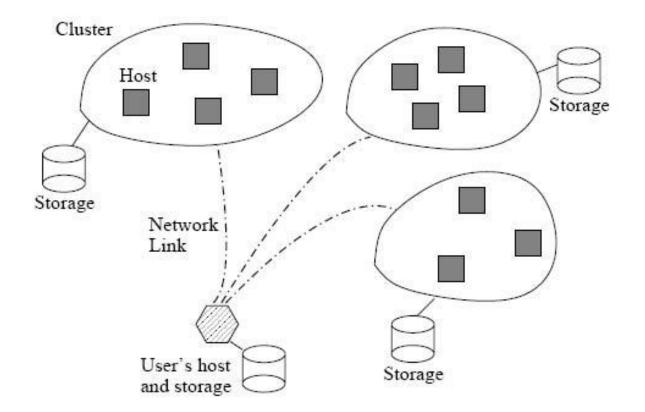


#### **Application Model**





#### System Model

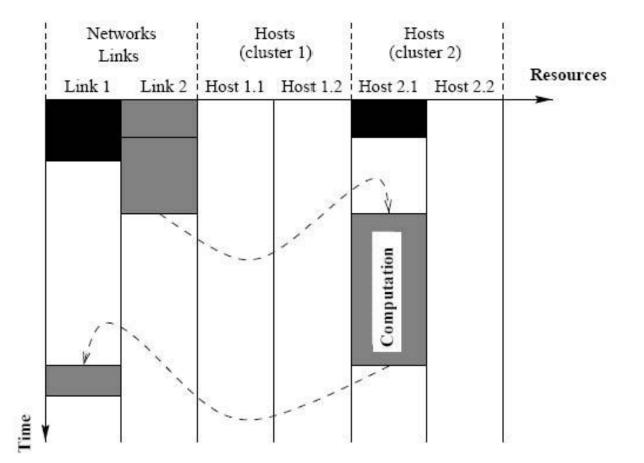


## **Scheduling Heuristic**

Schedule()

- 1. Compute the next Scheduling event
- 2. Create a Gantt Chart G
- 3. For each computation and file transfer underway
  - . Compute an estimate of its completion time
  - 2. Update the Gantt Chart G
- 4. Select a subset of tasks that have not started execution: T
- 5. Until each host has been assigned enough work
  - 1. Heuristically assign tasks to hosts
- 6. Convert G into a plan

#### Sample Gantt Chart



#### **Possible Variations**

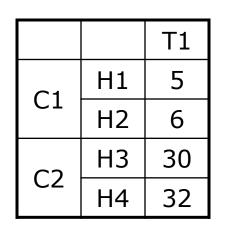
Schedule()

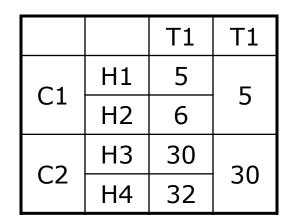
- Compute the next Scheduling event
- 2. Create a Gantt Chart G
  - For each computation and file transfer underway
    - . Compute an estimate of its completion time
    - Update the Gantt Chart G
- Select a subset of tasks that have not started execution: T
- Until each host has been assigned enough work
  - Heuristically assign tasks to hosts
- 6. Convert G into a plan



#### XSufferage

- Tasks may have little intra-cluster
   Completion time (CT) variation and large inter-cluster CT variation.
- Cluster-level MCT for Sufferage.







#### Simulations

- 1000 Simulated Grids, U(2,12)xU(2,32)
- 2000 applications, U(2,10)xU(20x1000)
- Task runtime U(100,300)
- Large Input File, U(400,100000) KBytes
- Small Input File, 1 KB
- One Output File, 10 KB
- Background load on the host machines and network links based on NWS traces
- Results over 1000 random Grid/application pairs.



	Geometric mean (sec)	Average Degradation from Best (%)	Average Rank
Max-min	2390	17.3	3.1
Min-min	2452	21.2	3.0
Sufferage	2329	14.1	2.8
XSufferage	2174	6.2	1.8

## Scheduling Task Graphs

- Task Graphs have dependencies between the tasks in the Application
- Scheduling methods for bag of task applications cannot be directly applied

Т5

5

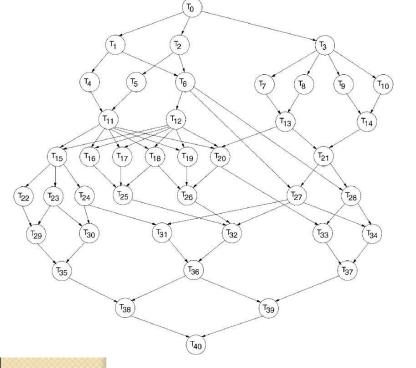
T<sub>13</sub>

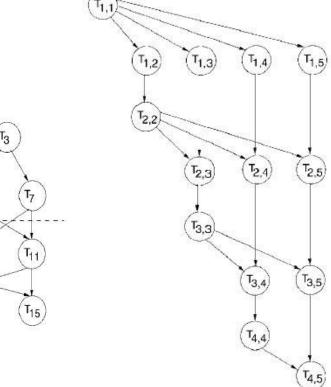
T12

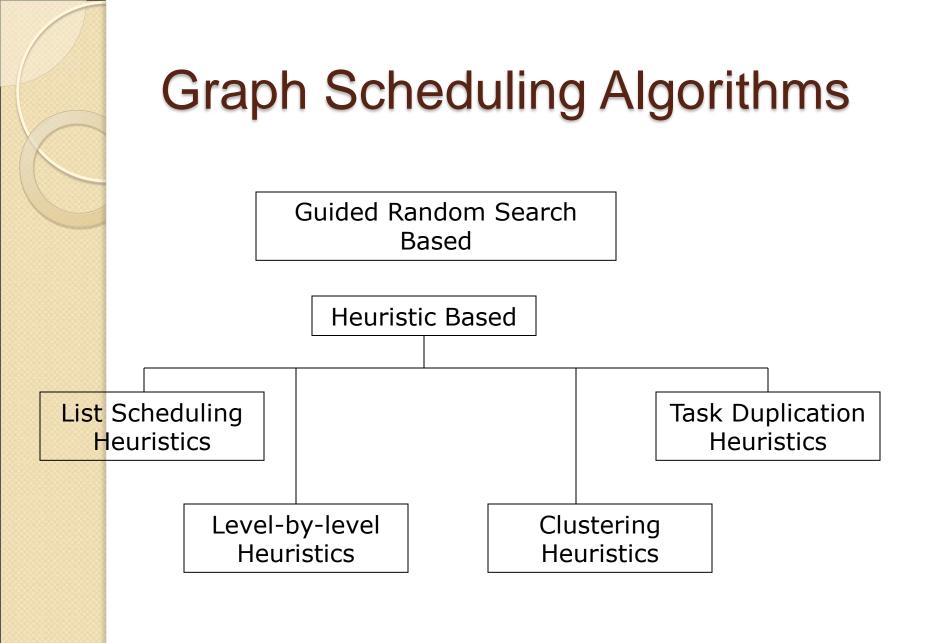
 $T_6$ 

T10

T14





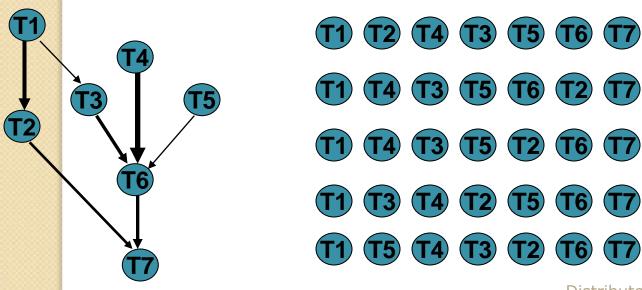


## Guided Random Search Based

- Genetic Algorithms
  - A chromosome is an ordering of tasks
  - A rule is required to convert it to a schedule
- Simulated Annealing
- Local Search Techniques, taboo etc

## **List Scheduling Heuristics**

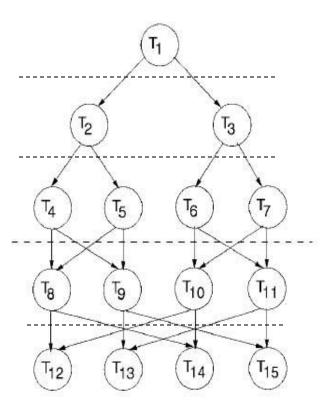
- An ordered list of tasks is constructed by assigning priority to each task
- Tasks are selected on priority order and scheduled in order to minimize a predefined cost function
- Tasks have to be in a topologically sorted order



#### Level by Level Scheduling

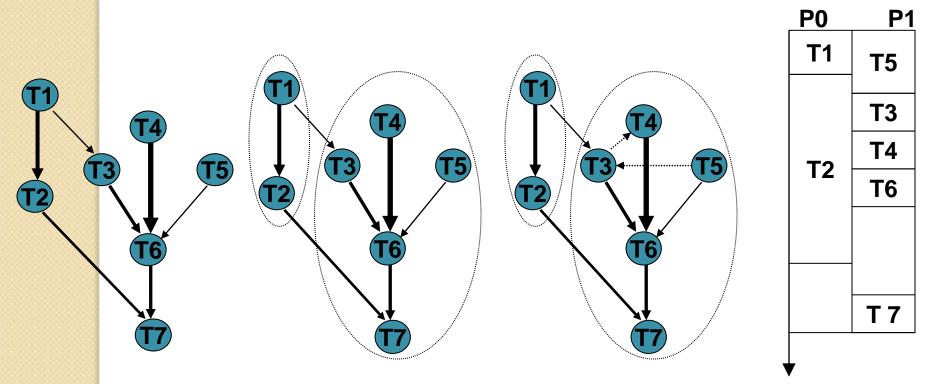
 Partition a DAG into multiple levels such that task in each level are independent.

 Apply Min-Min, Max-Min or other heuristics to tasks at each level



## **Clustering Heuristics**

Clustering heuristics cluster tasks together
Tasks in the same cluster are scheduled on the same processor



Distributed Operating

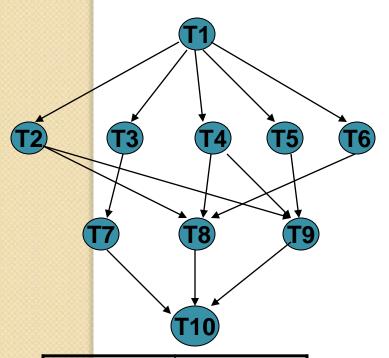
## Heterogeneous Earliest Finish Time

- List scheduling based heuristic
- Do a bottom up traversal of the graph and assign ranks to each task

$$rank_{u}(n_{i}) = \overline{w}_{i} + \max_{n_{j} \in succ(n_{i})}(\overline{c}_{i,j} + rank_{u}(n_{j}))$$
$$rank_{u}(n_{exit}) = \overline{w}_{exit}$$

## HEFT- contd

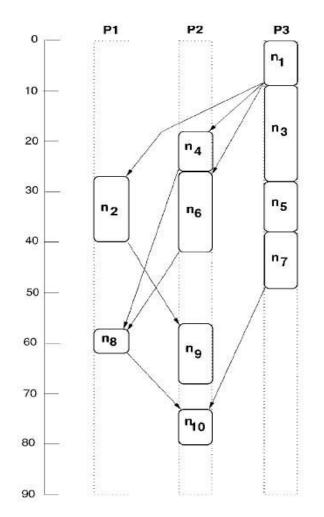
- Compute rank for all tasks in the graph
- Sort the tasks by non-increasing rank values (ensures topological sort)
- While there are unscheduled tasks
  - Select the first task in the list
  - Assign the task to the processor/machine that minimizes its completion time using insertion based scheduling
- endWhile



	Priority
T1	108
Т2	77
Т3	80
Т4	80
T5	69
Т6	63.33
T7	42.667
<b>T</b> 8	35.667
Т9	44.333
T10	14.667



**HEFT Schedule** 



## Critical Path on a Processor (CPOP)

Upward ranking

 $rank_u(n_i) = \overline{w}_i + \max_{n_j \in succ(n_i)} (\overline{c}_{i,j} + rank_u(n_j))$ 

• Downward ranking

 $rank_{d}(n_{i}) = \overline{w}_{i} + \max_{n_{j} \in pred(n_{i})}(\overline{c}_{i,j} + \overline{w}_{j} + rank_{d}(n_{j}))$ 

 $rank_d(n_{entry}) = 0$ 

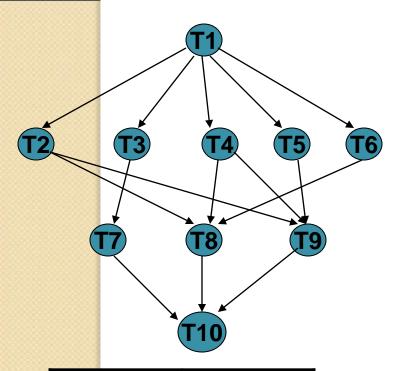
 $priority(n_i) = rank_u(n_i) + rank_d(n_i)$ 



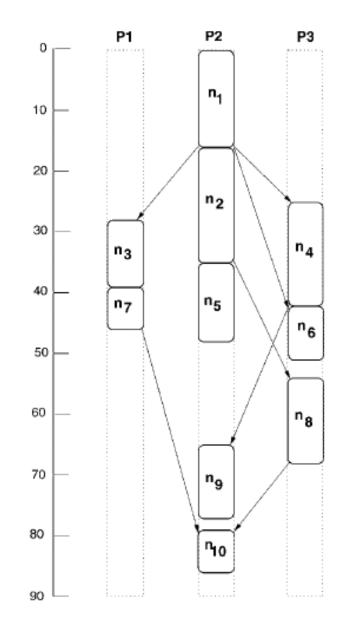
#### CPOP

$$\begin{split} |\mathsf{CP}| &= \mathsf{priority}(\mathsf{n}_{\mathsf{entry}}) \\ \mathsf{SET}_{\mathsf{CP}} &= \{\mathsf{n}_{\mathsf{entry}}\} \\ \mathsf{n}_k \leftarrow \mathsf{n}_{\mathsf{entry}} \\ \mathsf{While} \ \mathsf{n}_k \text{ is not the exit task do} \\ &\quad \mathsf{Select} \ \mathsf{n}_j \text{ where } ((\mathsf{n}_j \in \mathsf{succ}(\mathsf{n}_k) \text{ and } \mathsf{priority}(\mathsf{n}_j) == |\mathsf{CP}| \\ &\quad \mathsf{SET}_{\mathsf{CP}} = \{\mathsf{SET}_{\mathsf{CP}} \cup \mathsf{n}_j\} \\ &\quad \mathsf{n}_k \leftarrow \mathsf{n}_j \\ &\quad \mathsf{endWhile} \\ &\quad \mathsf{Select} \ \mathsf{the} \ \mathsf{Critical} \ \mathsf{Path} \ \mathsf{processor} \ \mathsf{p}_{\mathsf{cp}} \ \mathsf{that} \ \mathsf{minimizes} \ \mathsf{the} \ \mathsf{sum} \ \mathsf{of} \\ &\quad \mathsf{runtimes} \ \mathsf{of} \ \mathsf{tasks} \ \mathsf{on} \ \mathsf{the} \ \mathsf{critical} \ \mathsf{path} \end{split}$$

Go through the task list in priority order, assign tasks in  $\text{SET}_{\text{CP}}$  to  $p_{\text{cp}}$  and other tasks to any processor that minimizes its finish time



	Priority
T1	108
Т2	108
Т3	105
Т4	102
Т5	93
Т6	90.333
Т7	105
Т8	102.334
Т9	108
<b>T10</b>	108





#### Conclusions

- Heuristics for scheduling independent and dependent tasks on distributed systems
- Rescheduling in order to adapt to dynamic Grid conditions
- Partitioning in case of task graphs
- Future lecture (Nov 27<sup>th</sup>) on resource provisioning for applications.



#### ASSIGNMENT

 Q: Explain various scheduling algorithms in distributed system.